

Proof of Reachability, Part I.

$$\begin{aligned}
 X_0 &= 0 \\
 X_1 &= [R] \cup \text{Pr}_1(X_0) \\
 X_2 &= [R] \cup \text{Pr}_1(X_1)
 \end{aligned}$$

$$\begin{aligned}
 &\vdots \\
 &\downarrow \\
 X^* &= \lim_{n \rightarrow \infty} X_n.
 \end{aligned}$$

Define

π_0 arbitrary.

π_1 from $\text{Pr}_1(X_0)$

π_2 from $\text{Pr}_1(X_1)$

π_i is the move distr. that yields $\text{Pr}_1(X_{i-1})$.

Thm
For all $\epsilon > 0$, starting from S ,
 γ can reach R with prob $\geq X^*(S) - \epsilon$.

Note: If at s , γ play $\pi_i(s)$, the adversary plays η (a distribution of moves),

$$E_S^{\pi_i, \eta}(X_{i-1}) \geq \text{Pr}_1(X_{i-1}).$$

Here Pl. 2 may play suboptimally.

$$\text{Pr}_1(x_{i-1}) = \sup_{\xi} \inf_{\eta} E^{\xi\eta}(x_{i-1})$$

where

$$E^{\xi\eta}(x_{i-1}) = \sum_a \sum_b \sum_t p(s, a, t)(e) \cdot \xi(a) \cdot \eta(b).$$

Let ξ_i^* ^{← optimal for pr. 1.} = arg sup_ξ inf_η E^{ξη}(x_{i-1}).

Then,

$$E^{\xi_i^* \eta}(x_{i-1}) \geq \inf_{\eta} E^{\xi_i^* \eta}(x_{i-1})$$

$$= \sup_{\xi} \inf_{\eta} E^{\xi\eta}(x_{i-1}) = \text{Pr}_1(x_{i-1}) \quad \underline{\underline{\text{Q.E.D.}}}$$

Proof of Thm (Lemma 1 page 10):

You give me $\epsilon > 0$. I choose $n \in \mathbb{N}$ so that $X_n(s) \geq X^*(s) - \epsilon$.

Play $\pi_1^{(n)} = \xi_n \xi_{n-1} \xi_{n-2} \dots \xi_0 \xi_0 \xi_0 \dots$

Prove by induction on $n \geq 0$ that:

$$\forall \pi_2: \Pr_S^{\pi_1^{(n)}, \pi_2}(\text{OR}) \geq X_n(s).$$

Base case: $X_0 = 0$ obvious.

Ind step: s & R

$$\Pr_S^{\pi_1^{(n)}, \pi_2}(\text{OR}) \geq \sum_t \Pr_t^{\pi_1^{(n-1)}, \pi_2}(\text{OR}).$$

prob of OR from t
 prob of $s \rightarrow t$.

$\cdot \Pr_S^{\pi_1^{(n)}, \pi_2}$ (go to t) in 1 step

by ind hyp,

$$\geq \sum_t X_{n-1}(t) \cdot \Pr_S^{\pi_1^{(n)}, \pi_2}(\text{go to } t \text{ in 1 step})$$

$$\gg \sum_t X_{n-1}(t) \cdot Pr_s^{\xi_n \eta} \quad \begin{matrix} \text{(go to } t) \\ \text{in 1 step} \end{matrix} \quad \begin{matrix} \text{for some } \eta, \\ \eta = \pi_2(s). \end{matrix}$$

$$\gg E_s^{\xi_n \eta} (X_{n-1}) \quad \text{with } \dots$$

$$\gg Pr_{e_1} (X_{n-1})(s) = X_n(s) \quad \text{etc.}$$

$$X^* = \mu X. ([R] \sqcup Pre, (X))$$

$$\langle 1 \rangle \diamond R \gg \mu X. ([R] \sqcup Pre, (X)) .$$

↑
Consequence of the theorem.

Safety Games

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Goal: $\square B$.

$$X_0 = \mathbb{1}$$

$$X_1 = [B] \cap \text{Pre}, (X_0) = B \quad \text{Prds of } B \text{ for 1 step moment.}$$

$$X_2 = [B] \cap \text{Pre}, (X_1) \quad \text{max prds of } B \text{ for 2 moments.}$$

$$X_3 = [B] \cap \text{Pre}, (X_2) \quad \text{max prob of } B \text{ for 3 moments.}$$

⋮

$$X^* = \lim_{n \rightarrow \infty} X_n = \nu X. ([B] \cap \text{Pre}, (X)).$$

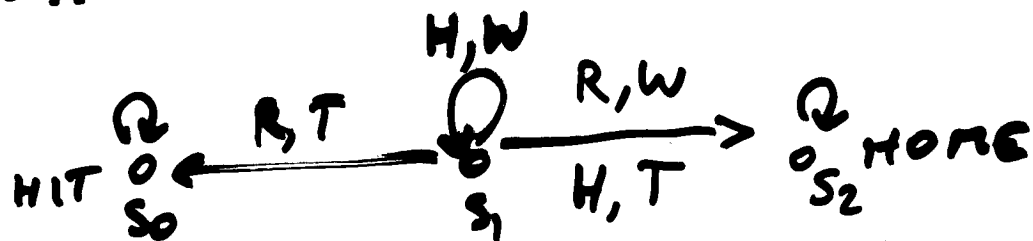
$$\langle 1 \rangle \square B \stackrel{=}{=} \nu X. ([B] \cap \text{Pre}, (X)).$$

$$\langle 1 \rangle \diamond B \stackrel{=}{=} \mu X. ([B] \cup \text{Pre}, (X)).$$

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Homework (next Monday):

Consider HIDE-OR-RUN, and compute



$$\Gamma_1(s_1) = \{R, W\}$$

$$\Gamma_2(s_1) = \{T, W\}$$

1) Compute $\chi_n(s_1)$ for $n \geq 0$.

2) Compute $\xi_n(s_1)$ for $n \geq 0$.